

Software Validation and Verification – 17/07/2026

Recall: Write your name on every paper sheet **and** motivate all your answers

Exercise 1

[6 points]

Assume that the set of propositions AP coincides with the natural numbers $\mathbb{N}$.

Let E_1 and E_2 be the linear time properties defined as follows:

$$E_1 = \left\{ \sigma = A_0, A_1, \dots \mid \bigcup_i A_i = \mathbb{N} \right\} \quad E_2 = \left\{ \sigma = A_0, A_1, \dots \mid \sum_i \sum_{n \in A_i} n = 110 \right\}$$

Informally, each word of E_1 contains all natural numbers, whereas each word of E_2 adds up to 110.

Express each of E_1 and E_2 as the intersection of a safety and of a liveness properties.

Exercise 2

[6 points]

Discuss whether the following linear time property is ω -regular:

$$E = \{ \sigma = A_0, A_1, \dots \mid \forall i, j. \text{ if } a \in A_i \text{ and } a \in A_j \text{ then } \#\{ k \mid i < k < j \text{ and } b \in A_k \} \text{ is even} \}$$

Informally, a word is in E if the occurrences of the proposition a are separated by an even number of states where b holds.

Notation: given a set X , we write $\#X$ for the *cardinality* of X , i.e. the number of its elements.

Exercise 3

[9 points]

Consider the set of LTL formulas $\Xi = \{\phi, \psi, \xi\}$, where:

$$\phi = \Box(a \Rightarrow (a \cup \bigcirc a)) \quad \psi = \Box(\neg a \Rightarrow \bigcirc a) \quad \xi = \neg \Box \Diamond \neg a$$

For each pair $(x, y) \in \Xi \times \Xi$, discuss whether $Word(x) \subseteq Word(y)$.

Exercise 4

[5 points]

Discuss the validity of the following statement:

*For any LTL formula ϕ , for any non-atomic proposition ψ subformula of ϕ ,
if $B' \xrightarrow{A} B$ and $B'' \xrightarrow{A} B$ are valid transitions of $\mathcal{G}_\phi$, then $\psi \in B'$ implies $\psi \in B''$*

Notation: given an LTL formula ϕ , we write $\mathcal{G}_\phi$ for the GNBA such that $\mathcal{L}_\omega(\mathcal{G}_\phi) = Word(\phi)$ returned by the algorithm from the lecture.

Exercise 5

[6 points]

Consider the following transition system and CTL formula.

$$\Phi = \forall \Diamond \exists ((\exists \Box a) \cup ((\exists \bigcirc b) \wedge (\exists \Box \neg a)))$$

Compute $Sat(\Phi)$ with the algorithm from the lecture.

